# **Approach for node-ID negotiation in CANopen**

The increasing complexity and size of CANopen (FD) networks create new challenges on how to decide which node-ID should be used. This article discusses a theoretical approach how devices in a CANopen (FD) network could negotiate their node-ID by themselves.

The assignment of a node-ID to a CANopen device is an essential requirement. Without a node-ID a CANopen device is unable to communicate over services defined in CiA 301 and CiA 1301. Without a valid node-ID in the range of 1 to 127, a CANopen device would never leave the NMT (network management) sub-state reset communication. Since the node-ID has to be unique in the network, there has to be a way to configure the node-ID value. Neither CiA 301 nor CiA 1301 specify a way to configure this value. There are various solutions on the market how this can be achieved.

Some manufactures use hardware solutions on their devices, such as dip- or rotary-switches, or encode plugs. Others, use some kind of proprietary software configuration. All of these solutions need more or less knowledge about the system where the devices are integrated. Another way to [configure the node-ID is specified in the document CiA 305,](https://can-newsletter.org/engineering/system-design/210607_canopen-fd-devices-identification-via-lss_cnlm_cia)  CANopen layer setting services (LSS). This specification describes services and protocols for identifying CANopen devices and assigning node-IDs, which can be used by a so called LSS manager (formerly named LSS master) or some configuration tool.

The CANopen application profile for building door control, CiA 416, also specifies a procedure for claiming node-IDs by the devices themselves. But besides that, this procedure is slow, it is also patented. The following theoretical approach for negotiation of node-IDs introduced by the engineers of Emotas Embedded Communication should be a starting point how to overcome most of the disadvantages of the other solutions.

# **Systems without a superordinated manager**

Some systems are designed very flexible. That means that a system could be built with different combinations of devices, exactly meeting the needed requirements.

Sometimes, it could be enough that only two devices are needed, sometimes much more, and sometimes there are multiple of the same device type in a network. An example of such kind of systems are heating systems based on CANopen. They can include multiple sensors and controllers depending on location and size where the system is installed.

When creating the devices for such a system, it is not clear which kind of device could be unique in the system, so it can operate as a CANopen manager. The CANopen manager is the device with NMT manager functionality plus additional functionality, for example, the LSS manager.

One way to assign node-IDs without an LSS manager is when integrating the system. But this leads to a static system, without a real plug-and-play possibility. And, in case some parts are added, or have to be replaced, an unused node-ID has to be assigned to the replaced part, which might be impossible for a service technician.

Another example are battery clusters, which are capable of plug and play. The only devices in this network are the batteries themselves. They are physically all of the same device type and have the same software. In such a system, the node-ID is irrelevant to the system, but is required for CANopen communication.

# **Systems with a superordinated manager**

In generic CANopen systems with a superordinated manager, there are also possibilities that the actual distribution of the node-IDs is not relevant for the functionality of the system.

Examples are the CANopen application profile for energy management systems, CiA 454, or the CANopen application profile for special-purpose car add-on devices, CiA 447. These specifications use the LSS Fastscan procedure, defined in CiA 305, to detect devices and to assign node-IDs to them. With LSS Fastscan devices can be only detected one after another, and in the worst case of completely unknown devices it will take 128 messages to verify every single bit of the CANopen LSS address. With a response-timeout of 10 ms, it would take 1,28 s to detect one single device.

# **New approach**

The idea behind this approach is that, in systems which do not depend on the node-ID distribution, the devices negotiate their node-ID themselves. The idea is not new, for example in the document J1939-81 of the SAE (Society of Automotive Engineers), there is a description of a so-called address claim procedure. In this procedure, the devices in a CANbased network negotiate the addresses (node-IDs), depending on values of the Name field. The problem here is that this procedure uses the 8-byte data field of the CAN frame. This led to the fact, that, in case two or more devices are claiming the same address, it could lead to collisions on the CAN network. In our approach, we are trying to avoid this by not using the data field: All the information exchanged between devices are encoded in the identifier field.

The main requirements for this procedure are that the software of the devices is able to send and receive CAN data frames in classical extended frame format (CEFF). Because most modern CAN FD controllers support sending and receiving this kind of data frames, this approach works in CANopen FD as well. It is also important that the software can distinguish between extended and basic frame formats.

#### **Usage of the CAN-Identifier**

The CAN-Identifier (CAN-ID) of an extended data frame has 29 bits available. For our approach we only need 13 bits. The least significant 13 bits are used for the node-ID negotiation. The 13 bits are divided into two fields, the least significant 8 bits are used as data field, and the most significant 5 bits as data code. Table 1 illustrates the usage of the CAN-ID.

*Table 1: 29-bit CAN-ID usage*

<b>Bit 28 to 13</b>	<b>Bit 12 to 8</b>	Bit 7 to 0
reserved	Data code	Data field

Using 13 bit of the 29 bits of an extended CAN-ID means that 44 % of the bits are required, but the following

$$
\frac{2^{13}}{2^{29}} = \frac{8.192}{536.870.912} = 0,00001525
$$

formula shows that only 0,0015 % of all possible extended CAN-IDs are reserved for this service:

The data code defines the meaning of the data field content. If the most significant bit is set to 1, the data field contains parts of the CANopen object  $1018<sub>h</sub>$ , identity object. For this procedure to work, all four sub-indices are required. Table 2 shows which part of the identity object  $1018<sub>h</sub>$  is transmitted with which data code.

Data code	1018 $b$ data	
10000 <sub>h</sub>	Sub-index 1, bit 0 to 7	
10001 <sub>b</sub>	Sub-index 1, bit 8 to 15	
10010 <sub>b</sub>	Sub-index 1, bit 16 to 23	
10011 <sub>b</sub>	Sub-index 1, bit 24 to 31	
10100 <sub>h</sub>	Sub-index 2, bit 0 to 7	
10101 <sub>b</sub>	Sub-index 2, bit 8 to 15	
10110 <sub>b</sub>	Sub-index 2, bit 16 to 23	
10111 <sub>b</sub>	Sub-index 2, bit 24 to 31	
11000 <sub>b</sub>	Sub-index 3, bit 0 to 7	
11001 <sub>b</sub>	Sub-index 3, bit 8 to 15	
11010 <sub>b</sub>	Sub-index 3, bit 16 to 23	
11011 <sub>b</sub>	Sub-index 3, bit 24 to 31	
11100 <sub>h</sub>	Sub-index 4, bit 0 to 7	
11101 <sub>b</sub>	Sub-index 4, bit 8 to 15	
11110 <sub>b</sub>	Sub-index 4, bit 16 to 23	
11111 <sub>b</sub>	Sub-index 4, bit 24 to 31	

*Table 2: Data code correlation with 1018h*

Due to the binary coding of the data code, it is very easy to implement the algorithm for accessing the value of the identity object. Bit 3 and 4 are the sub-index number - 1, and bit 0 and 1 are the byte number of this sub-index.

Two data code values are dedicated for the flow control of the process and are listed in Table 3.

*Table 3: Flow control data code*

Data code	<b>Name</b>
00001 <sub>b</sub>	RegUsedNodeld
00010 <sub>h</sub>	ActUsedNodeld

The data code ReqUsedNodeId is used to start a new process, whereby the data field contains 0. The resulting message with extended CAN ID  $100<sub>h</sub>$  will trigger all devices with a valid node-ID to respond with ActUsedNodeId and their own node-ID in the data field. Figure 1 shows the basic concept of asking the network for the node-IDs, which are already used in the network.

# **Negotiation process**

The negotiation process is only performed by nodes with an invalid node-ID. All nodes with a valid node-ID are ignoring the messages with the most significant bit in the data code field set to one. If a device participating in the procedure, receives a negotiation message, it compares the value according to Table 2 with its own equivalent. In case the received value is higher than its own value, the device continues the negotiation.

In case the received value is smaller as its own equivalent, the negotiation is stopped on this device. And, the device has to wait until the running negotiation is done. It can detect this by setting a timeout (e.g. 100 ms), and restart this timeout every time it receives a negotiation message. In case a device is able to transmit all 16 messages, which are needed to transmit all 128 bits of the identity object, it will then announce the preferred node-ID to the network by using ActUsedNodeId data code message.

Figure 2 shows a complete negotiation cycle with a timeout value of 100 ms before and after the actual negotiation and an additional 5-ms delay between every negotiation message.

# **Node-ID conflicts**

In case a device successfully negotiated a node-ID, but another device indicates by the ActUsedNodeId message the very same node-ID, the receiving device sends out the same ActUsedNodeId, sets its pending node-ID to the invalid node-ID, and enters the NMT state reset communication. Because the other device received this message, it loses its node-



*Figure 1: Initializing negotiation (Source: Emotas)*

ID too. Both devices can restart the negotiation process again. The losing device of this negotiation shall alter the preferred node-ID for the next cycle to a node-ID not yet  $\triangleright$ present in the system.

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*Figure 2: Complete negotiation cycle (Source: Emotas)*

emotas Virtual - emotas Virtual (500 kBit/s) - 0 % ●







Table 4 shows a CAN Trace simulation example of two nodes negotiating their node-IDs. The difference in the 1018h parameter revision number causes the one node to stop the negotiation and to restart again after the first one is ready.

The main goal is still to find a reliable, robust, and fast solution to assign node-IDs in plug-and-play systems to extend the network during its lifecycle. This theoretical approach should only be a basis for discussion on how to achieve this. With the discussed process it is possible to assign node-IDs to devices without the need of a dedicated LSS manager. The usage of extended CAN-IDs makes it possible that only 16 messages are needed, and that the process does not create CAN network collisions.

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